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## A Note on Anti-Berge Equilibrium for Bimatrix Game

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**Abstract.** We introduce a new concept of equilibrium based on Nash and Berge equilibriums. This equilibrium is called Anti-Berge equilibrium. We prove an existence of Anti-Berge equilibrium in the game. Based on Mills theorem [9], we reduce finding Anti-Berge equilibrium to a quadratic programming problem with linear constraints. The proposed approach has been illustrated on an example.

**Keywords:** Berge equilibrium, optimization, bimatrix game, Anti-Berge equilibrium.

### 1. Introduction

Game theory plays an important role in applied mathematics, economics and decision theory. There are many works devoted to game theory. Most of them deals with a Nash equilibrium. A global search algorithm for finding a Nash equilibrium was proposed in [13]. Also, the extraproximal and extragradient algorithms for the Nash equilibrium have been discussed in [3]. Berge equilibrium is a model of cooperation in social dilemmas, including the Prisoner's Dilemma games [15].

The Berge equilibrium concept was introduced by the French mathematician Claude Berge [5] for coalition games. The first research works of Berge equilibrium were conducted by Vaisman and Zhukovskiy [18; 19]. A method for constructing a Berge equilibrium which is Pareto-maximal with respect to all other Berge equilibriums has been examined in Zhukovskiy [10]. Also, the equilibrium was studied in [16] from a view point of differential games. Abalo and Kostreva [1; 2] proved the existence theorems for pure-strategy Berge equilibrium in strategic-form games of differential games. Nessah [11] and Larbani, Tazdait [12] provided with a new existence theorem. Applications of Berge equilibrium in social science have been discussed in [6; 17]. Also, the work [7] deals with an application of Berge equilibrium in economics. Connection of Nash and Berge equilibriums has been shown in [17]. Most recently, the Berge equilibrium was examined in Enkhbat and Batbileg [14] for Bimatrix game with its nonconvex optimization reduction. In this paper, inspired by Nash and Berge equilibriums, we introduce a new notion of equilibrium so-called Anti-Berge equilibrium. The main goal of this paper is to examine Anti-Berge equilibrium for bimatrix game.

The work is organized as follows. Section 2 is devoted to the existence of Anti-Berge equilibrium in a bimatrix game for mixed strategies. In Section 3, an optimization formulation of Anti-Berge equilibrium has been formulated.

## 2. Bimatrix Game

Consider the bimatrix game in mixed strategies with matrices  $(A, B)$  for players 1 and 2.

$$A = (a_{ij}), \quad i = 1, \dots, m,$$

$$B = (b_{ij}), \quad j = 1, \dots, n.$$

Denote by  $X$  and  $Y$  the sets

$$X = \{x \in R^m \mid \sum_{i=1}^m x_i = 1, x_i \geq 0, i = 1, \dots, m\},$$

$$Y = \{y \in R^n \mid \sum_{j=1}^n y_j = 1, y_j \geq 0, j = 1, \dots, n\}.$$

A mixed strategy for player 1 is a vector  $x = (x_1, x_2, \dots, x_m)^T \in X$  representing the probability that player 1 uses a strategy  $i$ . Similarly, the mixed strategies for player 2 is  $y = (y_1, y_2, \dots, y_n)^T \in Y$ . Their expected payoffs are given by :

$$f_1(x, y) = x^T Ay, \quad f_2(x, y) = x^T By.$$

First, we introduce the definitions of the equilibriums

**Definition 1.** A pair strategy  $(x^1, y^1) \in X \times Y$  is a Nash equilibrium if

$$\begin{cases} f_1(x^1, y^1) \geq f_1(x, y^1), & \forall x \in X, \\ f_2(x^1, y^1) \geq f_2(x^1, y), & \forall y \in Y. \end{cases}$$

**Definition 2.** A pair strategy  $(x^2, y^2) \in X \times Y$  is a Berge equilibrium if

$$\begin{cases} f_1(x^2, y^2) \geq f_1(x^2, y), & \forall y \in Y, \\ f_2(x^2, y^2) \geq f_2(x, y^2), & \forall x \in X. \end{cases}$$

**Definition 3.** A pair strategy  $(x^3, y^3) \in X \times Y$  is an Anti-Berge equilibrium (with respect to player 2) if

$$\begin{cases} f_1(x^3, y^3) \geq f_1(x^3, y), & \forall y \in Y, \\ f_2(x^3, y^3) \leq f_2(x, y^3), & \forall x \in X. \end{cases}$$

It is clear that

$$f_1(x^3, y^3) = \max_{y \in Y} f_1(x^3, y),$$

$$f_2(x^3, y^3) = \min_{x \in X} f_2(x, y^3).$$

**Definition 4.** A pair strategy  $(x^4, y^4) \in X \times Y$  is an Anti-Berge equilibrium (with respect to player 1) if

$$\begin{cases} f_1(x^4, y^4) \leq f_1(x^4, y), & \forall y \in Y, \\ f_2(x^4, y^4) \geq f_2(x, y^4), & \forall x \in X. \end{cases}$$

In Nash equilibrium both of players maximizes their payoff functions simultaneously. In Berge equilibrium both of players mutually supports each other to maximize their payoffs while in the Anti-Berge equilibrium one of them minimizes other's payoff function. In other words, one of them behaves unpleasantly and is antagonistic to other.

Before we introduce Anti-Berge equilibrium for 3-person game, it is worth mentioning Berge equilibrium [10] for the game.

**Definition 5.** A triple strategy  $(x^*, y^*, z^*) \in X \times Y \times Z$  is a Berge equilibrium if

$$\begin{cases} \hat{f}_1(x^*, y^*, z^*) \geq \hat{f}_1(x^*, y, z), & \forall (y, z) \in Y \times Z, \\ \hat{f}_2(x^*, y^*, z^*) \geq \hat{f}_2(x, y^*, z), & \forall (x, z) \in X \times Z, \\ \hat{f}_3(x^*, y^*, z^*) \geq \hat{f}_3(x, y, z^*), & \forall (x, y) \in X \times Y, \end{cases}$$

where the functions  $\hat{f}_i(x, y, z), i = 1, 2, 3$  defined on a set  $X \times Y \times Z$  of strategies are payoff functions of the players.

Now we introduce Anti-Berge equilibrium in the following.

**Definition 6.** A triple strategy  $(x^*, y^*, z^*) \in X \times Y \times Z$  is an Anti-Berge equilibrium (with respect to player 3) if

$$\begin{cases} \hat{f}_1(x^*, y^*, z^*) \geq \hat{f}_1(x^*, y, z), & \forall (y, z) \in Y \times Z, \\ \hat{f}_2(x^*, y^*, z^*) \geq \hat{f}_2(x, y^*, z), & \forall (x, z) \in X \times Z, \\ \hat{f}_3(x^*, y^*, z^*) \leq \hat{f}_3(x, y, z^*), & \forall (x, y) \in X \times Y. \end{cases}$$

An existence of Anti-Berge equilibrium for a bimatrix game is given by the following proposition.

**Theorem 1.** There exists an Anti-Berge equilibrium in a bimatrix game for mixed strategies.

*Proof.* We follow up similarly the proof done for Berge equilibrium in [14]. Define the sets  $S_1(x)$  and  $S_2(y)$  as follows:

$$S_1(\bar{x}) = \left\{ \bar{y} \in Y \mid f_1(\bar{x}, \bar{y}) = \max_{y \in Y} f_1(\bar{x}, y) \right\},$$

$$S_2(\bar{y}) = \left\{ \bar{x} \in X \mid f_2(\bar{x}, \bar{y}) = \min_{x \in X} f_2(x, \bar{y}) \right\}.$$

Since the functions  $f_1$  and  $f_2$  are continuous and the sets  $X$ ,  $Y$  are compact then there exist  $\max_{y \in Y} f_1(\bar{x}, y)$ ,  $\min_{x \in X} f_2(x, \bar{y})$ . Thus  $S_1(x) \neq \emptyset$  and  $S_2(y) \neq \emptyset$ .

Introduce the mapping  $\mathcal{K}$  in the following:

$$\mathcal{K}: X \times Y \rightarrow S_1 \times S_2.$$

□

It is clear that if  $(x^*, y^*)$  is Anti-Berge equilibrium then  $(x^*, y^*) \in \mathcal{K}(x^*, y^*)$ . We show that  $\mathcal{K}$  is convex compact. Indeed, for any  $(\tilde{x}, \tilde{y}) \in \mathcal{K}(\bar{x}, \bar{y})$  and  $(\hat{x}, \hat{y}) \in \mathcal{K}(\bar{x}, \bar{y})$  we have

$$\begin{aligned} f_1(\bar{x}, \tilde{y}) &= \max_{y \in Y} f_1(\bar{x}, y), \\ f_2(\tilde{x}, \bar{y}) &= \min_{x \in X} f_2(x, \bar{y}), \\ f_1(\bar{x}, \hat{y}) &= \max_{y \in Y} f_1(\bar{x}, y), \\ f_2(\hat{x}, \bar{y}) &= \min_{x \in X} f_2(x, \bar{y}). \end{aligned}$$

Since  $f_1$  and  $f_2$  are bilinear functions, for  $\alpha \in [0, 1]$  these equalities imply that

$$f_1(\bar{x}, \alpha \tilde{y} + (1 - \alpha) \hat{y}) = \max_{y \in Y} f_1(\bar{x}, y),$$

$$f_2(\alpha\tilde{x} + (1 - \alpha)\hat{y}, \bar{y}) = \min_{x \in X} f_2(x, \bar{y}),$$

which means that

$$(\alpha\tilde{x} + (1 - \alpha)\hat{y}, \alpha\tilde{y} + (1 - \alpha)\hat{y}) \in \mathcal{K}(\bar{x}, \bar{y}).$$

Thus  $\mathcal{K}$  is convex.

On the other hand,  $\max_{y \in Y} f_1(\bar{x}, y)$  and  $\min_{x \in X} f_2(x, \bar{y})$  are continuous functions on  $X \times Y$ , then  $\mathcal{K}$  is continuous mapping. Since  $X$  and  $Y$  are compact then by Tikhonov theorem [8]  $\mathcal{K}$  is also compact.

Therefore, conditions of fixed point theorem [4] are satisfied.

Hence, there exists  $(x^*, y^*)$  such that

$$(x^*, y^*) \in \mathcal{K}(x^*, y^*)$$

with  $x^* \in S_2(y^*)$  and  $y^* \in S_1(x^*)$ .

This means that

$$f_1(x^*, y^*) = \max_{y \in Y} f_1(x^*, y) \geq f_1(x^*, y), \quad \forall y \in Y,$$

$$f_2(x^*, y^*) = \min_{x \in X} f_2(x, y^*) \leq f_2(x, y^*), \quad \forall x \in X$$

which proves the assertion.

For further purpose, it is useful to formulate the following theorem.

**Theorem 2.** *A pair strategy  $(x^*, y^*)$  is an Anti-Berge equilibrium if and only if*

$$f_1(x^*, y^*) \geq [x^{*T} A]_j, \quad j = 1, 2, \dots, n, \quad (2.1)$$

$$f_2(x^*, y^*) \leq [B y^*]_i, \quad i = 1, 2, \dots, m. \quad (2.2)$$

*Proof. Necessity.* Assume that  $(x^*, y^*)$  is an Anti-Berge equilibrium. Then by Definition 3, we have

$$f_1(x^*, y^*) \geq x^{*T} A y, \quad \forall y \in Y, \quad (2.3)$$

$$f_2(x^*, y^*) \leq x^T B y^*, \quad \forall x \in X. \quad (2.4)$$

In the first inequality (2.3), successively choose  $y = (0, 0, \dots, 1, \dots, 0)$  with 1 in each of the  $n$  spots, in (2.4) choose  $x = (0, 0, \dots, 1, \dots, 0)$  with 1 in each of the  $m$  spots. We can easily see that

$$f_1(x^*, y^*) \geq [x^{*T} A]_j, \quad j = 1, \dots, n,$$

$$f_2(x^*, y^*) \leq [B y^*]_i, \quad i = 1, \dots, m.$$

**Sufficiency.** Suppose that for a pair  $(x^*, y^*) \in X \times Y$ , conditions (3.11) and (3.12) are satisfied. We choose  $x \in X$ ,  $y \in Y$  and multiply (3.11) by  $y_j$  and (3.12) by  $x_i$  respectively. We obtain

$$y_j f_1(x^*, y^*) \geq [x^{*T} A]_j y_j, \quad j = 1, 2, \dots, n.$$

Summing up these inequalities and taking into account that  $\sum_{j=1}^n y_j = 1$ , we get

$$f_1(x^*, y^*) = \left( \sum_{j=1}^n y_j \right) f_1(x^*, y^*) \geq \sum_{j=1}^n \sum_{i=1}^m a_{ij} x_i^* y_j = x^{*T} A y.$$

By analogy, we also have

$$f_2(x^*, y^*) = \left( \sum_{i=1}^m x_i \right) f_2(x^*, y^*) \leq \sum_{i=1}^m \sum_{j=1}^n b_{ij} x_i y_j^* = x^T B y^*.$$

Thus, we arrive at

$$\begin{aligned} f_1(x^*, y^*) &\geq f_1(x^*, y), \quad \forall y \in Y, \\ f_2(x^*, y^*) &\leq f_2(x, y^*), \quad \forall x \in X, \end{aligned}$$

concluding that  $(x^*, y^*)$  is an Anti-Berge equilibrium. The proof is complete.  $\square$

### 3. Quadratic Programming Formulation of Anti-Berge Equilibrium

**Theorem 3.** *A pair strategy  $(x^*, y^*)$  is an Anti-Berge equilibrium (with respect to player 2) for the bimatrix game if and only if there exist scalars  $(p^*, q^*)$  such that  $(x^*, y^*, p^*, q^*)$  is a solution to the following quadratic programming problem :*

$$\max_{(x, y, p, q)} F(x, y, p, q) = x^T (A - B) y - p + q \quad (3.1)$$

subject to :

$$[x^T A]_j \leq p, \quad j = 1, \dots, n, \quad (3.2)$$

$$[B y]_i \geq q, \quad i = 1, \dots, m, \quad (3.3)$$

$$\sum_{i=1}^m x_i = 1, \quad x_i \geq 0, \quad i = 1, \dots, m, \quad (3.4)$$

$$\sum_{j=1}^n y_j = 1, \quad y_j \geq 0, \quad j = 1, \dots, n. \quad (3.5)$$

Proof can be done similarly to the theorem in [14] proven for a Berge equilibrium.

*Proof. Necessity.* Suppose that  $(x^*, y^*)$  is an Anti-Berge equilibrium. Choose scalars  $p^*$  and  $q^*$  such that  $p^* = f_1(x^*, y^*)$ ,  $q^* = f_2(x^*, y^*)$ .

We show that  $(x^*, y^*, p^*, q^*)$  is a solution to problem (3.1)–(3.5). First, we show that  $(x^*, y^*, p^*, q^*)$  is a feasible point for problem (3.1)–(3.5).

By Theorem 2, the equivalent characterization of an Anti-Berge equilibrium point, we have

$$p^* = f_1(x^*, y^*) \geq \left[ x^{*T} A \right]_j, \quad j = 1, \dots, n,$$

$$q^* = f_2(x^*, y^*) \leq [B y^*]_i, \quad i = 1, \dots, m.$$

The rest of the constraints are satisfied because of  $x^* \in X$  and  $y^* \in Y$ . It means that  $(x^*, y^*, p^*, q^*)$  is a feasible point.

Choose any  $x \in X$  and  $y \in Y$ . Multiply (3.2)–(3.3) by  $y_j$  and  $x_i$ , respectively. If we sum up these inequalities, we obtain

$$f_1(x, y) = x^T A y \leq p,$$

$$f_2(x, y) = x^T B y \geq q.$$

Hence, we get

$$F(x, y, p, q) = x^T (A - B) y - p + q \leq 0$$

for all  $x \in X$ ,  $y \in Y$ . But with  $p^* = f_1(x^*, y^*)$  and  $q^* = f_2(x^*, y^*)$ , we have  $F(x^*, y^*, p^*, q^*) = 0$ . Hence, the point  $(x^*, y^*, p^*, q^*)$  is a solution to problem (3.1)–(3.5).

**Sufficiency.** Let  $(\bar{x}, \bar{y}, \bar{p}, \bar{q})$  be a solution to problem (3.1)–(3.5).

We show that  $(\bar{x}, \bar{y})$  is an Anti-Berge equilibrium of the game. Since  $(\bar{x}, \bar{y}, \bar{p}, \bar{q})$  is a feasible point, the following constraints are satisfied:

$$[\bar{x}^T A]_j \leq \bar{p}, \quad j = 1, \dots, n, \quad \sum_{i=1}^m \bar{x}_i = 1, \quad \bar{x}_i \geq 0, \quad i = 1, \dots, m, \quad (3.6)$$

$$[B \bar{y}]_i \geq \bar{q}, \quad i = 1, \dots, m, \quad \sum_{j=1}^n \bar{y}_j = 1, \quad \bar{y}_j \geq 0, \quad j = 1, \dots, n, \quad (3.7)$$

Hence, we have

$$f_1(\bar{x}, \bar{y}) = \bar{x}^T A \bar{y} \leq \bar{p} \sum_{j=1}^n \bar{y}_j = \bar{p}, \quad (3.8)$$

$$f_2(\bar{x}, \bar{y}) = \bar{x}^T B \bar{y} \geq \bar{q} \sum_{i=1}^m \bar{x}_i = \bar{q}. \quad (3.9)$$

Summing up these inequalities, we obtain

$$F(\bar{x}, \bar{y}, \bar{p}, \bar{q}) = \bar{x}^T(A - B)\bar{y} - \bar{p} + \bar{q} \leq 0. \quad (3.10)$$

Taking into account (3.8) and (3.9), we conclude that the function  $F(x, y, p, q)$  reaches its maximum at zero:

$$F(\bar{x}, \bar{y}, \bar{p}, \bar{q}) = (\bar{x}^T A \bar{y} - \bar{p}) + (\bar{x}^T B \bar{y} - \bar{q}) = 0 \quad (3.11)$$

with

$$\bar{x}^T A \bar{y} = \bar{p}, \quad (3.12)$$

$$\bar{x}^T B \bar{y} = \bar{q}. \quad (3.13)$$

From (3.12)-(3.13) and (6)-(7) we have

$$\bar{p} = f_1(\bar{x}, \bar{y}) = \bar{x}^T A \bar{y} \geq [\bar{x}^T A]_j \quad j = 1, \dots, n,$$

$$\bar{q} = f_2(\bar{x}, \bar{y}) = \bar{x}^T B \bar{y} \leq [B \bar{y}]_i \quad i = 1, \dots, m.$$

Now by Theorem 2,  $(\bar{x}, \bar{y})$  is an Anti-Berge equilibrium which completes the proof.  $\square$

Note that the condition

$$F(x^*, y^*, p^*, q^*) = 0$$

is necessary and sufficient for a  $(x^*, y^*)$  to be an Anti-Berge equilibrium.

We can also formulate the following assertion for Anti-Berge equilibrium (with respect to player 1).

**Theorem 4.** *A pair strategy  $(\hat{x}^*, \hat{y}^*)$  is an Anti-Berge equilibrium (with respect to player 1) for the bimatrix game if and only if there exist scalars  $(\hat{p}^*, \hat{q}^*)$  such that  $(\hat{x}^*, \hat{y}^*, \hat{p}^*, \hat{q}^*)$  is a solution to the following quadratic programming problem :*

$$\max_{(x, y, p, q)} F(x, y, p, q) = x^T(B - A)y + p - q$$

subject to :

$$[x^T A]_j \geq p, \quad j = 1, \dots, n,$$

$$[B y]_i \leq q, \quad i = 1, \dots, m,$$

$$\sum_{i=1}^m x_i = 1, \quad x_i \geq 0, \quad i = 1, \dots, m,$$

$$\sum_{j=1}^n y_j = 1, \quad y_j \geq 0, \quad j = 1, \dots, n.$$



As an example, consider the following bimatrix game with matrices  $A$  and  $B$  :

$$A = \begin{pmatrix} 9 & 11 & 6 & 20 \\ 7 & 4 & 10 & 21 \\ 2 & 16 & 15 & 9 \\ 5 & 9 & 9 & 17 \\ 4 & 3 & 5 & 2 \end{pmatrix} \text{ and } B = \begin{pmatrix} 15 & 10 & 5 & 19 \\ 13 & 18 & 1 & 16 \\ 11 & 17 & 18 & 12 \\ 6 & 11 & 3 & 10 \\ 8 & 12 & 8 & 7 \end{pmatrix}$$

Problem (3.1)–(3.5) for finding Anti-Berge equilibrium (with respect to player 2) is formulated as:

$$\max F(x, y, p, q) = -6x_1y_1 + x_1y_2 + x_1y_3 + x_1y_4 - 6x_2y_1 - 14x_2y_2 + 9x_2y_3 + 5x_2y_4 - 9x_3y_1 - x_3y_2 - 3x_3y_3 - 3x_3y_4 - x_4y_1 - 2x_4y_2 + 6x_4y_3 + 7x_4y_4 - 4x_5y_1 - 9x_5y_2 - 3x_5y_3 - 5x_5y_4 + p - q$$

$$\left\{ \begin{array}{ll} 9x_1 + 7x_2 + 2x_3 + 5x_4 + 4x_5 - p & \leq 0 \\ 11x_1 + 4x_2 + 16x_3 + 9x_4 + 3x_5 - p & \leq 0 \\ 6x_1 + 10x_2 + 15x_3 + 9x_4 + 5x_5 - p & \leq 0 \\ 20x_1 + 21x_2 + 9x_3 + 17x_4 + 2x_5 - p & \leq 0 \\ x_1 + x_2 + x_3 + x_4 + x_5 = 1 & \\ 15y_1 + 10y_2 + 5y_3 + 19y_4 - q & \geq 0 \\ 13y_1 + 18y_2 + y_3 + 16y_4 - q & \geq 0 \\ 11y_1 + 17y_2 + 18y_3 + 12y_4 - q & \geq 0 \\ 6y_1 + 11y_2 + 3y_3 + 10y_4 - q & \geq 0 \\ 8y_1 + 12y_2 + 8y_3 + 7y_4 - q & \geq 0 \\ y_1 + y_2 + y_3 + y_4 = 1 & \\ x_1 \geq 0, x_2 \geq 0, x_3 \geq 0, x_4 \geq 0, x_5 \geq 0, & \\ y_1 \geq 0, y_2 \geq 0, y_3 \geq 0, y_4 \geq 0, y_5 \geq 0. & \end{array} \right.$$

We can easily check that  $F(x^*, y^*, p^*, q^*) = 0$  with  $x^* = (0, 0, 0, 0.273, 0.727)^T$ ,  $y^* = (0, 0, 0.375, 0.625)^T$ ,  $p^* = 6.09$ ,  $q^* = 7.375$  and  $F^* = 0$ . It means that  $(x^*, y^*)$  is an Anti-Berge equilibrium (with respect to player 2) for the bimatrix game.

On the other hand, the game has also Anti-Berge equilibrium (with respect to player 1) in pure strategies:  $x^* = (0, 1, 0, 0, 0)^T$ ,  $y^* = (0, 1, 0, 0)^T$ . But there are two another Anti-Berge equilibria:

$$x^1 = (0.8125, 0, 0.1875, 0, 0)^T, y^1 = (0.764706, 0.235294, 0, 0)^T,$$

$$x^2 = (0.532895, 0.447368, 0.019737, 0, 0)^T, y^2 = (0.6875, 0.21875, 0.09375, 0)^T.$$

## Conclusion

We examined so-called Anti-Berge equilibrium in a bimatrix game. By analogy of Nash and Berge equilibriums, we proved the existence of Anti-Berge equilibrium in the game. Finding an Anti-Berge equilibrium in the game has been reduced to a quadratic programming problem with an indefinite matrix. An example has been considered. We introduced also Anti-Berge equilibrium, a new concept of equilibria, for 3-person game. Computational aspects of Anti-Berge equilibria will be discussed in a next paper.

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## Равновесие анти-Бержа для биматричных игр

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**Аннотация.** Рассматривается новая биматричная игра на основе равновесий Нэша и Бержа. Решение данной игры будем называть равновесием анти-Бержа. С помощью теоремы Милса [9] задача нахождения равновесия анти-Бержа сводится к задаче квадратичного программирования с линейными ограничениями. Новое понятие равновесия анти-Бержа иллюстрируется на численном примере.

**Ключевые слова:** равновесие Бержа, оптимизация, биматричная игра, равновесие анти-Бержа.

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